

Re: Convex Hull of Points on a Straight Line

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If we take two points (x_1, y_1) and (x_2, y_2) then by a convex combination of these two points we mean any point of the following form:

$$[kx_1 + (1-k)x_2, ky_1 + (1-k)y_2] \text{ for } 0 \leq k \leq 1$$

In this context I would just like to add a point that in the lower bound proof where sorting is reduced to convex hull problem to avoid the problem of colinearity the data x for sorting is mapped to the point (x, x^2) on a parabola and not to (x, x) on a line for the 2D convex hull problem.

Thanks.

----- Pinaki

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eKo1 wrote:

What do you mean by "convex combination?" Would you provide an example?

Babua wrote:

This is a degenerate 2D problem. If you consider it as a 1D problem then any convex hull will be a line segment. Any convex combination of any two points within that segment also belongs to that segment. So the 1D problem can be solved in $O(n)$ time unlike the 2D problem that is lower bounded by $\omega(n \log n)$.

Thanks.

---- Pinaki

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eKo1 wrote:

Re: Convex Hull of Points on a Straight Line

Let L be a straight line on the cartesian plane. Pick a finite set of points S on that line. What is the convex hull of S ?

From my understanding of hull points, the convex hull is the two points

farthest away from each other in S . Is this true?

I ask because there is a version of Graham's algorithm in my discrete mathematics book that does not consider the case when there are two points whose segments with the first hull point have the same angle with respect to the horizontal. I guess this version of the algorithm is flawed in that respect.